

Derived Satake via factorization II

Recall from last time: we want to extend

$$\mathrm{Sph}_{\tilde{G}, \times}^{\mathrm{spec}} = \mathrm{IndCoh}(\mathcal{H}_{\tilde{G}, \times}^{\mathrm{spec}}) \cong \mathrm{Sym}(\check{\mathfrak{g}}[-2])\text{-mod}(R\mathrm{p}(\check{G}))$$

to a factorization category, and to lift
 B-F to an equivalence of monoidal
 fact categories $\mathrm{Sph}_G \cong \mathrm{Sph}_{\check{G}}^{\mathrm{spec}}$ ($\mathrm{Sph}_G := D(\mathcal{H}_G)^{\mathrm{ren}}$).

Heuristic: $f: Y \rightarrow Z$ affine morphism

$$Q(\mathrm{Coh}(Y(D_{\mathrm{dR}}))_{Z(D_{\mathrm{dR}})}^{\times} Z(D_{\mathrm{dR}})) \cong f_* \mathrm{Dy}\text{-mod}^{\mathrm{fact}}(Q(\mathrm{Coh}(Z(D_{\mathrm{dR}})))$$

Theorem (Raskin) For reasonable Z (e.g. classifying stacks), this is literally true at a point.

Aside (Betti setting): $A = \text{comm. alg.}$

$$A\text{-mod}^{\mathbb{H}\mathbf{E}_n} \xrightarrow{\mathrm{oblv}} \mathrm{Vect}$$

admits a left adjoint.

Proposition The monad

$$\text{Vect} \xrightarrow[\text{oblv}]{} A\text{-mod}^{IE_n}$$

is given the assoc alg. $S^{n-1} \otimes A$.

$$\Rightarrow A\text{-mod}^{IE_n} \xrightarrow{\sim} S^{n-1} \otimes A\text{-mod}^*$$

Note that $\text{Spec}(S^{n-1} \otimes A) = \underline{\text{Map}}(S^{n-1}, \text{Spec } A)$.

For $n=2$ we get a Beilis version of the heuristic:

$$A\text{-mod}^{IE_2} \xrightarrow{\sim} \mathcal{Q}\mathcal{G}\mathcal{H}(\underline{\text{Map}}(S^1, \text{Spec } A))$$

Our case: $\mathcal{Y} = pt/G$, $\mathcal{Z} = pt/G \times pt/G$, $\mathcal{A} = \Delta$

$$\mathcal{Y}(\overset{\circ}{D}_{dR})_{\mathcal{Z}(\overset{\circ}{D}_{dR})} \times \mathcal{Z}(\overset{\circ}{D}_{dR}) = LS_G(\overset{\circ}{D})_{LS_G(\overset{\circ}{D}) \times LS_G(\overset{\circ}{D})} \times LS_G(\overset{\circ}{D})_{LS_G(\overset{\circ}{D}) \times LS_G(\overset{\circ}{D})}$$

$$= LS_G(\overset{\circ}{D})_{LS_G(\overset{\circ}{D})} LS_G(\overset{\circ}{D}) = \mathcal{H}_G^{\text{Spec}}$$

$$\Rightarrow \mathcal{Q}\mathcal{G}\mathcal{H}(\mathcal{H}_{G, \times}^{\text{Spec}}) \cong \mathcal{O}_G\text{-mod}^{\text{fact}}(\text{Rep}(G \times G))_x$$

$I = \text{nonempty finite set}$

$$\rightsquigarrow \mathcal{O}_G\text{-mod}^{\text{tac}}(\text{Rep}(G \times G))_{X^I}$$

$\text{Rep}(G \times G)_{x_j}$

$$J = I \amalg J_0, \quad x_J = \{x_j\} \in X^J \rightsquigarrow M_{x_J} \in \text{Vect}$$

comp. as J varies, w/ iso.'s

$$M_{x_J} \cong M_{x_I} \otimes \bigotimes_{j \in J_0} \mathcal{O}_G$$

when $x_j \neq x_k$ for $j \neq k \in J_0$, and $x_i \neq x_j$ for $i \in I, j \in J_0$

Renormalization

Why renormalize?

- to get IndCoh (as opposed to Qcoh) at a point, needed for B-F;
- to get a better category, e.g. compactly generated
- to get a "ULA generated" category: once the derived Salatah functor is constructed, can check that it is an equivalence at a point.

\mathcal{T} = DG cat. w/ t-structure

Standing assumptions:

- $\ell^{\geq 0}$ stable under filtered colimits;
- (right completeness) $\operatorname{colim}_n \mathcal{T}^{\leq n} c \xrightarrow{\sim} c$ for any $c \in \mathcal{T}$.

An object $c \in \mathcal{T}$ is called almost compact if $\tau^{\geq n} c$ is compact in $\mathcal{T}^{\leq n}$ for all n .

Call c coherent if it is almost compact and eventually coconnective (almost compact \Rightarrow eventually connective by right completeness).

E.g. $\mathcal{T} = \mathcal{Q}\mathcal{G}\mathcal{h}(Y)$, Y = almost finite type scheme
 $f \in \mathcal{Q}\mathcal{G}\mathcal{h}(Y)$ is coherent in this sense
 \Leftrightarrow coherent in usual sense.

We say that the t-structure on \mathcal{T} is coherent if (in addition to standing assumptions)

- $\ell^{\geq 0}$ is compactly generated;
- $\ell^{\geq 1} \hookrightarrow \ell^{\geq 0}$ preserves compact objects
(\Leftrightarrow coherent objects preserved by truncation)

Remark: \mathcal{C} coherent $\Rightarrow (\mathcal{C}^\heartsuit)^c \subset \mathcal{C}^\heartsuit$ is abelian

Definitions are from appendix to "Integral
transforms for coherent sheaves" by Ben-Zvi,
Nadler, and Preygel.

If \mathcal{C} is coherent and is compactly generated
by $\mathcal{C}^{\text{coh}} :=$ coherent objects, we call \mathcal{C}
regular.

NB: coherence is a property of \mathcal{C}^+ (not so
for regularity).

If \mathcal{C} is coherent, then it admits a
canonical renormalization \mathcal{C}^{ren} t -exact

$$\mathcal{C}^{\text{ren}} := \text{Ind}(\mathcal{C}^{\text{coh}}) \xrightarrow{\sim} \mathcal{C}$$

such that

$$(\mathcal{C}^{\text{ren}})^+ \xrightarrow{\sim} \mathcal{C}^+.$$

E.g. $\mathcal{C} = \mathcal{O}(\text{coh}(\gamma))$, γ almost finite type
 \mathcal{C} coherent and $\mathcal{O}(\text{coh}(\gamma))^{\text{ren}} = \text{Ind}(\text{coh}(\gamma))$.