

# A Sixteenth-order Polylogarithm Ladder

Henri Cohen, Leonard Lewin and Don Zagier

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Using the LLL algorithm and the second author's "ladder" method, we find (conjectural)  $\mathbf{Z}$ -linear relations among polylogarithms of order up to 16 evaluated at powers of a single algebraic number. These relations are in accordance with theoretical predictions and are valid to an accuracy of 300 decimal digits, but we cannot prove them rigorously.

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## 1. INTRODUCTION

Let  $m$  be a positive integer. The  $m$ -th polylogarithm function is defined for complex  $z$  with  $|z| < 1$  by

$$\operatorname{Li}_m(z) = \sum_{n=1}^{\infty} \frac{z^n}{n^m}.$$

For  $m = 1$ ,  $\operatorname{Li}_1(z) = -\log(1 - z)$ , while for  $m > 1$ , the function  $\operatorname{Li}_m(z)$  is a higher transcendental function. The functions  $\operatorname{Li}_m(z)$  can be extended analytically to the whole complex plane if one makes a cut, for example, the real line from 1 to infinity (corresponding, for  $m = 1$ , to the principal branch of the logarithm).

The function  $\operatorname{Li}_1(z)$  satisfies the functional equation

$$\operatorname{Li}_1(1 - xy) = \operatorname{Li}_1(1 - x) + \operatorname{Li}_1(1 - y),$$

which implies that any linear combination with integral coefficients of  $\operatorname{Li}_1$  values of elements of a number field will again be an  $\operatorname{Li}_1$  value of an element of this number field. The corresponding statement for higher polylogarithms is not true. However, the  $\operatorname{Li}_m$  satisfy "trivial" functional equations relating  $\operatorname{Li}_m(x)$  to  $\sum_{y^p=x} \operatorname{Li}_m(y)$  for all integers  $p \neq 0$ ; and for small orders  $m$ , they also satisfy more complicated nontrivial functional equations. (The present record, due to H. Gangl, gives nontrivial functional equations for  $m = 7$ .) In addition, these functional equations give relations among polylogarithm values of elements of a given

number field. It is conjectured that *all* relations among polylogarithms of algebraic arguments can be obtained by specializing functional equations, but this is not known for any  $m \geq 2$ .

The main goal of this paper is to provide an example of “ladder” relations (i.e., relations among polylogarithms of powers of a fixed number  $\alpha$ ) for the number  $\alpha \cong 1.17628$  defined as the unique root outside the unit circle of the self-reciprocal polynomial

$$X^{10} + X^9 - X^7 - X^6 - X^5 - X^4 - X^3 + X + 1.$$

Dilogarithm relations for this number were first investigated by G. Ray and are given in Chapter 7 of [Lewin 1991]. The ladder relations we will find go up to order 16, confirming the general theory (as described in [Zagier 1991a] and recalled in Section 3). This is quite possibly the highest order occurring for any algebraic number, a property related to the well-known conjecture of Lehmer [1933] that  $\alpha$  has the smallest possible Mahler measure for an algebraic integer that is not a root of unity (the Mahler measure of an algebraic integer is equal to the product of the moduli of the number’s conjugates that are greater than 1, hence here equal to  $\alpha$ ). The computation of the ladder relations requires both high-accuracy computation of polylogarithms (we needed more than 300 decimal digits) and careful use of linear algebra and lattice reduction techniques for fairly large matrices with nonexact entries, so both of these issues will have to be addressed in the paper.

## 2. NUMERICAL COMPUTATION OF POLYLOGARITHMS

To achieve our goal, it is necessary to compute polylogarithms to a reasonably high accuracy. To do this, we use the following very simple and pretty formula for the polylogarithm, which we could not find in the literature:

**Proposition 1.** *For  $m$  a natural number and  $|z| < 2\pi$  we have*

$$\begin{aligned} \text{Li}_m(e^z) = & \sum_{\substack{n \geq 0 \\ n \neq m-1}} \zeta(m-n) \frac{z^n}{n!} \\ & + \left(1 + \frac{1}{2} + \cdots + \frac{1}{m-1} - \log(-z)\right) \frac{z^{m-1}}{(m-1)!}. \end{aligned}$$

The proof by induction is straightforward. We observe, as an amusing remark, that the polylogarithm of complex order  $s$  (defined for  $|z| < 1$  like  $\text{Li}_m$  but with  $m$  replaced by  $s$ ) can be analytically extended by the formula

$$\text{Li}_s(e^z) = \sum_{n \geq 0} \zeta(s-n) \frac{z^n}{n!} + \Gamma(1-s)(-z)^{s-1}$$

for all  $s \notin \mathbf{N}$ ; the limit of the right-hand side as  $s \rightarrow m \in \mathbf{N}$  is the expression occurring in the proposition.

If  $x = \rho e^{i\theta} \in \mathbf{C}$  and we write  $x = e^z$ , the condition  $|z| < 2\pi$  is equivalent to

$$\exp(-\sqrt{4\pi^2 - \theta^2}) < \rho < \exp(\sqrt{4\pi^2 - \theta^2}).$$

Since we may assume that  $|\theta| \leq \pi$ , this region contains the annulus  $e^{-\pi\sqrt{3}} < \rho < e^{\pi\sqrt{3}}$  (note that  $e^{-\pi\sqrt{3}} < 0.005$  and  $e^{\pi\sqrt{3}} > 230$ ). So, to compute  $\text{Li}_m(x)$ , we use the power series if  $|x|$  is small (say  $|x| \leq \frac{1}{2}$ ), the formula of proposition 1 if  $|x|$  is near 1 (say  $\frac{1}{2} < |x| < 2$ ), and the functional equation relating  $\text{Li}_m(x)$  to  $\text{Li}_m(x^{-1})$  if  $|x|$  is large (say  $|x| > 2$ ; in fact, as we will see, we will never have to use this).

There is still one more point to treat. We will need polylogarithm values to several hundred decimals, hence we also need the values of  $\zeta(k)$  for integers  $k$  to such a high accuracy. When  $k \leq 0$  or  $k > 0$  and even,  $\zeta(k)$  is given explicitly in terms of Bernoulli numbers  $B_n$ . We must also compute  $\zeta(k)$  for  $k$  odd,  $k \geq 3$ . We could use the standard Euler–McLaurin method for doing this, but there is a simpler and much faster method, coming from the theory of modular forms (more precisely, by integrating  $k$  times the Fourier expansion of the holomorphic Eisenstein series of weight  $k+1$ ); these formulas were known to Ramanujan.

**Proposition 2.** *For  $k$  odd,  $k > 1$ , we have*

$$\begin{aligned} \zeta(k) = & \frac{(2\pi)^k}{k-1} \sum_{n=0}^{(k+1)/2} (-1)^n (1-2n) \beta_{2n} \beta_{k+1-2n} \\ & - 2 \sum_{n \geq 1} \frac{n^{-k}}{e^{2\pi n} - 1} \left(1 + \frac{2\pi n \varepsilon_k}{1 - e^{-2\pi n}}\right), \end{aligned}$$

where  $\beta_n$  denotes  $B_n/n!$  and  $\varepsilon_k = 0$  for  $k \equiv 3 \pmod{4}$ ,  $\varepsilon_k = 2/(k-1)$  for  $k \equiv 1 \pmod{4}$ .

The above propositions give us rapidly convergent methods for computing polylogarithms in the

whole complex plane, and in particular near or on the unit circle. (Of course, Proposition 2 is not really necessary, since the values of  $\zeta(k)$  for  $k$  up to 16 could be computed just once by a slower method and then stored.)

We will actually be working with a modified version of the polylogarithm function which is defined by

$$P_m(x) = \operatorname{Re}_m \left( \sum_{r=0}^m \frac{2^r B_r}{r!} (\log|x|)^r \operatorname{Li}_{m-r}(x) \right),$$

where  $\operatorname{Re}_m$  denotes real or imaginary part depending on whether  $m$  is odd or even, and  $\operatorname{Li}_0$  has been set equal to  $-\frac{1}{2}$  [Zagier 1991a]. In contrast with the original polylogarithms, these modified functions satisfy “clean” functional equations (i.e., equations involving only polylogarithms of a given order, without lower-order correction terms). In particular, we have  $P_m(1/x) = (-1)^{m-1} P_m(x)$ .

### 3. LADDERS

In [Abouzahra and Lewin 1985] and [Lewin 1984], recipes to obtain linear relations among polylogarithm values of powers of a single algebraic integer  $\alpha$  are given. These recipes, which are special cases of the theory describing linear relations among polylogarithm values of arbitrary algebraic arguments (as explained in [Zagier 1991a; 1991b]), are as follows.

Assume that one knows, for some algebraic number  $\alpha$ , a certain number of relations of the form

$$\prod_{n=1}^{\infty} (\alpha^n - 1)^{c_n} = \zeta \alpha^N$$

for some integers  $c_1, c_2, \dots$  ( $c_n = 0$  for almost all  $n$ ) and  $N$  and some root of unity  $\zeta$ . Such relations are called *cyclotomic relations*. For the special number defined in the introduction, since  $\alpha$  is real and greater than 1, we must have  $\zeta = 1$ ; moreover, since  $1/\alpha$  is a conjugate of  $\alpha$ , we deduce by conjugating the equation that  $N$  must be equal to  $\frac{1}{2} \sum_n n c_n$ . Since the modified unilogarithm function  $P_1$  is given by  $P_1(x) = -\log|1-x| + \frac{1}{2} \log|x|$ , the cyclotomic equation can be rewritten

$$\sum_{n=1}^{\infty} c_n P_1(\alpha^n) = 0.$$

So, these are our unilogarithm relations, the first rung of the ladder. The main statement of the theory is that all relations among  $m$ -th polylogarithms are obtained by  $(m-1)$ -fold *pseudointegration*, that is, the formal replacement of  $P_1(\alpha^n)$  by  $n^{-(m-1)} P_m(\alpha^n)$ . In other words, any ladder relation at the  $m$ -th rung should have the form

$$\sum_{n=1}^{\infty} \frac{c_n}{n^{m-1}} P_m(\alpha^n) = 0$$

for some cyclotomic relation  $\prod (\alpha^n - 1)^{c_n} = \alpha^N$ .

However, not all relations obtained by pseudointegration will hold. The rest of the picture is as follows. First, observe that the set of cyclotomic equations forms a group under multiplication, so the set of possible collections of exponents  $\{c_n\}$  forms a group under addition; obviously, the same is true for the linear relations among the polylogarithms of any given order. Suppose that we have found  $I(1)$  multiplicatively independent cyclotomic relations. Then, starting with  $m = 2$ , suppose inductively that we have  $I(m-1)$  linearly independent equations at order  $m-1$ , say

$$S_{m-1,i}(\alpha) = \sum_{n=1}^{\infty} \frac{c_n^{(i)}}{n^{m-2}} P_{m-1}(\alpha^n) = 0$$

for  $i = 1, \dots, I(m-1)$ . We suppose further that the same relation remains true if  $\alpha$  is replaced by any of its conjugates  $\alpha^\sigma$ . (In practice, and according to conjectures, this is automatically true, but in any case, we need it for the inductive setup.) We write this symbolically as  $S_{m-1,i}(\alpha^\sigma) = 0$ . We then consider the pseudointegrated quantities

$$S_{m,i}(\alpha^\sigma) = \sum_{n=1}^{\infty} \frac{c_n^{(i)}}{n^{m-1}} P_m((\alpha^\sigma)^n).$$

There are  $r_1 + 2r_2$  possible embeddings  $\sigma$  of  $\mathbf{Q}[\alpha]$  into the complex numbers, where  $r_1$  and  $r_2$  denote as usual the number of real and complex conjugates of  $\alpha$ , so for each  $i$  the collection of numbers  $\{S_{m,i}(\alpha^\sigma)\}_\sigma$  can be considered as a vector in  $\mathbf{R}^{r_1+2r_2}$ . However, because of the functional equations of the function  $P_m(x)$ , not all of the components of this vector are independent, so we can think of it as a vector in a smaller dimensional vector space  $\mathbf{R}^{d(m)}$  of dimension  $d(m)$  depending on the parity of  $m$ . Specifically, since  $P_m(x)$  is invariant under  $x \mapsto \bar{x}$  for  $m$  odd and anti-invariant for

$m$  even, we need only half of the complex conjugates and can ignore the real conjugates when  $m$  is even, reducing the number of components from  $r_1 + 2r_2$  to  $r_1 + r_2$  or  $r_2$  (in our special case, 6 and 4), respectively. Moreover, for  $\alpha$  like ours satisfying a self-reciprocal equation, the functional equation

$$P_m(1/x) = (-1)^{m-1} P_m(x)$$

implies that we also need only half of the real embeddings when  $m$  is odd, i.e.,  $d(m)$  is  $r_2$  (in our case, 4) when  $m$  is even, and  $r_1/2 + r_2$  (in our case, 5) when  $m$  is odd. Now the theory says that the vectors obtained by pseudointegrating the valid  $(m-1)$ -th order relations will lie in a lattice in  $\mathbf{R}^{d(m)}$ , or in other words, that any  $d(m) + 1$  of them will satisfy a linear relation over  $\mathbf{Z}$  rather than merely over  $\mathbf{R}$ . (This assertion, given conjecturally in [Zagier 1991a], is actually now a theorem, to appear in forthcoming work by Beilinson and Deligne.) In particular, at least  $I(m) = I(m-1) - d(m)$  of our  $I(m-1)$  vectors in  $\mathbf{R}^{d(m)}$  will be 0. Changing bases if necessary, we can suppose that these are simply the relations  $S_{m,i}(\alpha^\sigma) = 0$  for  $i = 1, \dots, I(m)$ , and this puts us back in the same situation as before and lets us continue the inductive process. Summarizing, in climbing the ladder, we lose (at most)  $d(m)$  relations at each step and hence, starting with the  $I(1)$  cyclotomic relations, can mount the ladder as long as the quantity  $I(m) = I(1) - d(2) - \dots - d(m)$  remains positive. Note that, although the theory tells us that such relations exist, it does not give any bound on the size of the coefficients. As a consequence, although the relations which we will give are correct to several hundred decimal places, and although the number and type of relations we find agrees with the number which we know must exist, the specific relations that we give are not rigorously proved.

A striking example of a “ladder” of the sort just described is given in [Abouzahra et al. 1987] and [Zagier 1991a], where one takes for  $\alpha$  a root of  $X^3 - X - 1 = 0$ . In this case, one can climb the ladder up to level 8, and at level 9, there remain 2 pseudointegrated relations which are just sufficient to give a relation between nonalgorithms if one adds  $P_9(\alpha^0) = P_9(1) = \zeta(9)$  to the list of admissible powers of  $\alpha$ . In the present paper, we study the special Salem number mentioned in the introduction, an even more spectacular example.

The first step is to get the ladder started by finding cyclotomic relations for the number  $\alpha$  as explained at the beginning of this section. We do this as follows. First notice that the polynomial  $X^n - 1$  factors as a product of cyclotomic polynomials  $\Phi_k(X)$  over all divisors  $k$  of  $n$ , so that finding multiplicative relations among  $\alpha$  and the numbers  $\alpha^n - 1$  is the same as finding multiplicative relations among  $\alpha$  and the numbers  $\Phi_k(\alpha)$ . Any such relation will imply a relation among the norms, so we first compute the norm of  $\Phi_k(\alpha)$  for many  $k$ . We find that this norm is  $\pm 1$  for 66 values of  $k$  less than 1000, namely,

$k = 1, 2, 3, 5, 6, 7, 8, 9, 10, 11, 13, 16, 17, 18, 20, 21, 23, 24, 27, 28, 29, 30, 34, 37, 38, 40, 42, 44, 45, 47, 50, 52, 56, 57, 60, 63, 64, 65, 66, 70, 74, 75, 76, 78, 84, 86, 92, 96, 98, 105, 110, 118, 132, 138, 144, 154, 160, 165, 186, 195, 204, 212, 240, 270, 286, 360$ .

But even when the norm of  $\Phi_k(\alpha)$  is not a unit, we can hope to (multiplicatively) combine several  $\Phi_k(\alpha)$  to obtain a unit, and this indeed happens, giving us the nine further units  $\Phi_{k_1}(\alpha)/\Phi_{k_2}(\alpha)$  for  $(k_1, k_2) = (12, 4), (36, 4), (62, 31), (108, 4), (124, 31), (130, 26), (175, 25), (182, 14), (246, 82)$ .

Together with  $\alpha$  itself, this gives us 76 multiplicatively independent units and hence, since the unit rank of  $K = \mathbf{Q}[\alpha]$  is equal to  $r_1 + r_2 - 1 = 5$ , 71 linearly independent cyclotomic relations. (The five units  $\alpha$  and  $\Phi_k(\alpha)$  for  $k = 1, 2, 3, 5$  form a generating system of independent units.)

Using Baker-type methods, it is possible to give an effective upper bound on the index of the cyclotomic polynomials which have to be examined. However, as often with this type of bound, it does not seem possible to decrease it to a computationally feasible value. Hence, although it is possible that some cyclotomic relation has escaped our search, this seems quite unlikely.

#### 4. CLIMBING THE LADDER

By the construction sketched above, from our initial  $I(1) = 71$  cyclotomic relations, we must obtain  $I(2) = 67$  dilogarithm relations,  $I(3) = 62$  trilogarithm relations, decreasing alternatively by 4 or 5, and finally we must obtain  $I(16) = 4$  hexadecalogarithm relations. (We would need six such

relations to be able to climb to the 17th level.) To obtain these relations is not a completely trivial task, and in this section we explain the method that has been used to obtain them. All the computations were done using the GP/PARI calculator developed by C. Batut, D. Bernardi, H. Cohen and M. Olivier.

There are exactly 111 different exponents of  $\alpha$  which occur in the cyclotomic relations. Hence, we will represent the cyclotomic relations as a  $111 \times 71$  matrix  $L_1$  with integral entries, each column representing a relation. This matrix is, of course, not unique, and we can try to get the coefficients as small as possible by using the LLL algorithm [Lenstra et al. 1982]. We can give a rough measure of the size of the coefficients by computing the  $L^2$ -norm of  $L_1$ . In the present case, the minimal value found (and probably the absolute minimum) is 20.

Now we pseudointegrate these relations, which simply means that we divide each of the 111 rows  $j$  by its corresponding exponent (from 1 to 360), which we denote by  $e(j)$ . Let  $M_2$  be the resulting matrix. We form the  $4 \times 111$  matrix  $V_2 = (v_{i,j})$  with  $v_{i,j} = P_2(\alpha_i^{e(j)})$ , where the  $\alpha_i$  are four nonreal conjugates of  $\alpha$  (one for each complex conjugate pair), and then form the  $4 \times 71$  matrix  $V_2 \cdot M_2$  of 4-tuples of dilogarithm values for the 71 pseudointegrated relations. This matrix turns out to have maximal rank (4), so we pick four independent columns (actually, the first four), and then form the  $71 \times 67$  matrix

$$K_2 = \ker(V_2 \cdot M_2),$$

whose 67 columns are the relations obtained by writing each of the remaining (last) 67 columns of  $V_2 \cdot M_2$  as a linear combination of the chosen four. (Notice that this procedure is numerically well-defined and stable even though the entries of our matrices are only approximate, because the invertibility of the  $4 \times 4$  matrix depends on the nonvanishing of a number, and this is something which can be checked numerically.) The theory described in Section 3 tells us that  $K_2$  should, in fact, have rational coefficients, and it is then easy to see that a possible choice for the matrix  $L_2$  of relations for dilogarithms would then be  $L_2 = M_2 \cdot K_2$ .

The first problem that faces us is that  $K_2$  has only been calculated approximately, so we have to somehow recognize its entries as rational numbers.

Because  $M_2$  has integer entries and we are only concerned with the product of this matrix with  $K_2$ , we may as well first compute

$$L'_2 = M_2 \cdot K_2$$

and then try to approximate it by a matrix with rational entries. If we are given sufficient accuracy, we can simply use the continued fraction expansion of the coefficients and hope to find a reasonably simple denominator. In the present case, the coefficients are apparently all integers, so we have nothing more to do than rounding, but in higher levels, we will have to find a common denominator and define  $L_m$  by multiplying  $L'_m$  by this denominator and rounding. (Even here, if we had not LLL-reduced the matrix  $L_1$ , we would have found the rather small denominator 6.)

If  $L_2$  is our rounded matrix, we can say with confidence (especially if we work with 300 decimal digits, which we will need for higher levels) that we have found the correct 67 relations at the dilogarithm level. This is however still not entirely satisfactory, since we may not have a  $\mathbf{Z}$ -basis of the lattice of relations, but only a  $\mathbf{Q}$ -basis.

To check whether we have a  $\mathbf{Z}$ -basis, we use the following easily proved lemma.

**Lemma.** *Let  $M$  be an  $m \times n$  matrix with integer entries and rank  $n$ . Then the lattice spanned by the columns of  $M$  is equal to the intersection with  $\mathbf{Z}^m$  of the  $\mathbf{Q}$ -vector space spanned by the columns of  $M$  if and only if the GCD of all  $n \times n$  subdeterminants is equal to 1.*

We will say that a matrix satisfying the condition of the preceding lemma is *primitive*. To replace a matrix  $M$  by a primitive one whose columns generate the same  $\mathbf{Q}$ -vector space, we can, for instance, proceed as follows. We extract at random two  $n \times n$  determinants and let  $d$  be their GCD. Then, for every prime  $p$  dividing  $d$ , we make  $M$  primitive at  $p$  by elementary linear algebra modulo  $p$ . Note that for this, we do not need to factor  $d$  completely, since once the small prime factors have been removed, we can treat the remaining part as a prime, and the only thing which can go wrong is that we obtain a nontrivial factor of this “prime,” in which case we try again with that factor and its cofactor. Alternatively, we could apply Hermite reduction to  $M$ , after which it is obvious how to make it primitive.

Applying this procedure to  $L_2$ , we now obtain a  $\mathbf{Z}$ -basis of the space of relations, not just a  $\mathbf{Q}$ -basis. As before, we terminate that matrix by LLL-reducing. In our case, if we start from the LLL-reduced matrix  $L_1$ , the GCD  $d$  of subdeterminants is already equal to 1, so the reduction step is unnecessary. On the other hand, if we had started from a non-LLL-reduced matrix  $L_1$  giving a denominator 6, we would have had to eliminate 6 (i.e., 2 and 3) from the GCD of the subdeterminants.

Since we LLL-reduce each matrix  $L_m$  as we proceed, the relations that we obtain are essentially optimal. Finding a kernel is no problem, since it involves inverting a  $4 \times 4$  or  $5 \times 5$  matrix followed by a matrix multiplication. Similarly, the process of making a matrix primitive is fairly straightforward. The crucial problem which arises now is that of explosion of denominators. Since the kernel is computed in an arbitrary way at first, and only corrected afterwards, it is not surprising that we have little control on the denominators which occur. In fact, working still with 300 decimals of accuracy, we are able to go not much further than level 9 or 10. So we must still improve our strategy.

A first improvement is as follows. It is clear that the pseudointegrated matrices  $M_m$  (after clearing denominators) may not be primitive in the above sense. In fact, the GCD of maximal subdeterminants will be divisible by all the primes dividing the exponents of the cyclotomic relations. In our case, this corresponds exactly to all the primes up to and including 59. We can make the matrices  $M_m$  primitive with respect to all of these primes, and experiment shows that no other primes need to be eliminated. With this improvement, we were able to reach level 12. However, once again, we were caught up by the explosion of the size of the denominators, i.e., after computing  $L'_{13} = M_{13} \cdot \ker(V_{13} \cdot M_{13})$ , we were unable to recognize the common denominator of the coefficients of  $L'_{13}$ .

At that point, several ideas were tried. First, we could simply have increased the precision of the polylogarithm values, but a rough estimate showed that we probably would have exceeded reasonable computer storage and time limitations. Second, since the coefficients of  $L'_{13}$  must have a *common* denominator, it might be possible to use the LLL algorithm to find it, by using a version which is essentially dual to the version used in finding linear

relations between real numbers. However, we did not succeed with that either. The idea which finally worked is very simple: Before computing the kernel, do an LLL-reduction of the matrix  $M_m$ . It is not unreasonable to expect that this will lower not only the size of the coefficients of  $M_m$ , but also that of the denominators. This is, in fact, strikingly true (see next section), and enabled us to climb to the top of the ladder (i.e.,  $m = 16$ ) in a few more minutes and thus to obtain the desired four hexadecalogarithm relations between powers of  $\alpha$ .

### Remarks

(1) It would be of considerable interest to develop an algorithm which can compute directly an LLL-reduced  $\mathbf{Z}$ -basis of relations between linearly dependent vectors  $v_1, \dots, v_n$  without computing first a  $\mathbf{Q}$ -basis of  $\sum \mathbf{Q}v_i$ , then transforming this into a  $\mathbf{Z}$ -basis, and finally LLL-reducing. The algorithm should be able to deal with exact or imprecise entries, i.e., we should not have to first find an  $\mathbf{R}$ -basis and then “recognize” the real entries as rational numbers. One way to do this is to apply LLL-reduction to the lattice  $\mathbf{Z}^n$  equipped with the quadratic form

$$Q(a_1, \dots, a_n) = \|a_1v_1 + \dots + a_nv_n\|^2 + \varepsilon(a_1^2 + \dots + a_n^2)$$

for  $a_1, \dots, a_n \in \mathbf{Z}$ , with a suitably small  $\varepsilon$ ; then the very short vectors correspond to the kernel and automatically give an LLL-reduced basis of it. (When the Euclidean space in which the  $v_i$  lie is one-dimensional, this is the standard way of finding integer relations among real numbers.) This algorithm can be made to work but requires careful choice of  $\varepsilon$ .

(2) Initially, instead of using the conjugates of  $\alpha$ , we simply looked (using LLL) for relations among the unmodified polylogarithm functions  $\text{Li}_m$  evaluated at powers of the real number  $\alpha$ . However, the size of the denominators and of the coefficients which occur increases even more rapidly than in the method we finally used, and we were unable to climb beyond level 7.

## 5. NUMERICAL RESULTS

Let us summarize the procedure described above. Let  $L_1$  be the  $111 \times 71$  LLL-reduced matrix of cy-

clotomic relations. Consider some integer  $m > 1$  and assume inductively that the matrix  $L_{m-1}$  has been computed. Let  $M_m$  be the matrix obtained after performing successively the following steps:

- (a) pseudointegrate the matrix  $L_{m-1}$ ;
- (b) reduce the columns to integer entries having GCD equal to 1;
- (c) eliminate the primes up to 59;
- (d) LLL-reduce.

Let  $V_m$  be the  $d(m) \times 111$  ( $d(m) = 4$  or  $5$ ) matrix of polylogarithm values of the powers of the relevant conjugates of  $\alpha$ . Set  $L'_m = M_m \cdot \ker(V_m \cdot M_m)$ , where the kernel is computed by simple matrix inversion. Using the continued fraction expansion of a few coefficients of  $L'_m$ , determine a plausible denominator  $d_m$  for  $L'_m$ , and let  $L_m$  be the matrix obtained from  $L'_m$  by the following steps:

- (a) Multiply  $L'_m$  by  $d_m$  and then round each coefficient to the nearest integer; let  $\varepsilon_m$  be the sum of the absolute values of the distances between each coefficient and its rounded value.
- (b) Eliminate  $d_m$  from the resulting matrix. In practice, eliminate individually the small prime factors of  $d_m$ , and globally the remaining cofactor, even if it is not prime.
- (c) LLL-reduce.

In Table 1, we give the results obtained by working with 305 decimals of accuracy. Note that since there is some freedom in the LLL-reduction procedure, another implementation may give slightly different results, although, of course, the lattices of polylogarithm relations will be the same. We give successively  $m$ ,  $\varepsilon_m$ ,  $d_m$  and  $l_m$ , the  $L^2$  norm of the LLL-reduced matrix  $L_m$ .

A few remarks are in order concerning this table. First, thanks to the combination of all the tricks mentioned above, the explosion of denominators has been completely avoided for small levels, and in fact, the denominator is equal to 1 for all even levels up to 10. On the other hand, when the level is 13 or above, the denominators become very large. We believe that the reason for this is that the matrices  $M_m$  get narrower as  $m$  grows (for example,  $M_{16}$  has only eight columns), hence the LLL-reduction of these matrices which saved us in the end is unfortunately not of much help. A related observation is that up to level 12, the precision of the results stays excellent, but thereafter

$m$	$\varepsilon_m$	$d_m$	$l_m$
1	—	—	$2 \cdot 10^1$
2	$10^{-305}$	1	$4.4 \cdot 10^1$
3	$10^{-304}$	2	$1.4 \cdot 10^2$
4	$10^{-305}$	1	$6.4 \cdot 10^2$
5	$10^{-301}$	24	$8.8 \cdot 10^3$
6	$10^{-305}$	1	$2.2 \cdot 10^4$
7	$10^{-297}$	81	$1.0 \cdot 10^6$
8	$10^{-305}$	1	$2.2 \cdot 10^6$
9	$10^{-294}$	63057	$2.3 \cdot 10^7$
10	$10^{-305}$	1	$6.6 \cdot 10^8$
11	$10^{-286}$	234280024	$1.3 \cdot 10^{11}$
12	$10^{-291}$	12	$6.9 \cdot 10^{13}$
13	$10^{-265}$	$2.5 \cdot 10^{24}$	$1.5 \cdot 10^{18}$
14	$10^{-259}$	$1.0 \cdot 10^{25}$	$7.4 \cdot 10^{23}$
15	$10^{-204}$	$4.6 \cdot 10^{72}$	$7.4 \cdot 10^{37}$
16	$10^{-153}$	$1.1 \cdot 10^{106}$	$2.7 \cdot 10^{71}$

TABLE 1. Salem number ladder data

we suffer a loss of precision of about the same order as  $d_m l_m$ . In particular, if we had worked with only 250 decimals instead of 300, we would not have been able to obtain the hexadecalogarithm relations. By using the method sketched earlier in Remark (1) preceding this section, we could have avoided the denominator explosion altogether, but apparently the same precision, about 300 digits, would still have been needed to obtain the highest-order relation.

The third remark is that the relations have been LLL-reduced, hence the  $L^2$ -norms given in the last column are certainly close to best possible, up to maybe a variation of 1 or 2 in the exponent of 10. In particular, the four relations that we obtain at level 16 have 70-digit coefficients, and the LLL bounds tell us that no such relation exists whose coefficients have at most 69 digits. Just for fun, we have given one of these relations in the Appendix.

The last remark is the following. At even orders  $m$ , we have worked with four nonreal conjugates of  $\alpha$ , since this is enough to obtain the relations, and also because the modified polylogarithm function  $P_m$  vanishes at real arguments. However, the theory tells us that for each relation  $\sum c_n P_m((\alpha^\sigma)^n) = 0$ , the corresponding com-

bination  $\sum c_n L_m(\alpha^n)$  of the modified *real* polylogarithm function

$$L_m(x) = \sum_{r=0}^{m-2} \frac{(-\log|x|)^r}{r!} \text{Li}_{m-r}(x) - \frac{(-\log|x|)^{m-1}}{m(m-2)!} \log|1-x|$$

for  $m$  even and  $-1 \leq x \leq 1$  should give a rational multiple of  $\pi^m$ . This is indeed what we find. Moreover, the calculation lets us answer a question that had arisen in earlier polylogarithm investigations but could not previously be answered because the orders of the polylogarithms occurring were too small. In all known relations among the values of the real polylogarithm function  $L_m$  at real arguments, the combination of polylogarithms which occur are not merely rational multiples of  $\pi^m$ , but are rational multiples with a small denominator. Conjecturally, a prime  $p$  can occur in the denominator only if some extension of  $F$  of degree  $m$  contains a  $p$ -th root of unity; in particular, for generic fields, not having especially large intersections with

cyclotomic fields, one expects (at most) the same denominator as for  $\mathbf{Q}$ , i.e., at most  $D_m$ , if we write  $\zeta(m) = \zeta_{\mathbf{Q}}(m)$  as  $N_m \pi^m / D_m$  with  $N_m$  and  $D_m$  coprime integers. For instance, all known relations among values of the values of the function  $L_2$  (“Rogers dilogarithm function”) at arguments in real quadratic fields other than  $\mathbf{Q}(\sqrt{5})$  give integral multiples of  $\pi^2/6 = \zeta(2)$ . The question was whether the right general conjecture should be that polylogarithm relations give simple rational multiples of  $\pi^m$ , or simple rational multiples of  $\zeta(m)$ , i.e., whether for generic  $F$ , they are integral multiples of  $N_m \pi^m / D_m$  or merely of  $\pi^m / D_m$ . This could not be checked before because no relations beyond order 9 were known, and the first  $m$  with  $N_m > 1$  is 12 ( $N_{12} = 691$ ). It turned out that for  $m = 12$  and  $m = 16$ , the rational multiples occurring did *not* have numerators divisible by  $N_m$ . For instance, the combination of modified hexadecalogarithms given in the Appendix is an integer multiple of  $\pi^{16}/D_{16} = \pi^{16}/325641566250$ , but not of  $\zeta(16) = 3617\pi^{16}/325641566250$ .

**APPENDIX: A (CONJECTURAL) HEXADECALOGARITHM RELATION**

As promised, we give one of the four relations among the hexadecalogarithm values of powers of  $\alpha$ . We give the coefficients in tabular form, each line containing an integer  $n$  between 1 and 360, and a coefficient  $c_n$  of at most 71 digits such that the sum  $\sum_n c_n P_{16}((\alpha^\sigma)^n)$  vanishes for each conjugate  $\alpha^\sigma$  of  $\alpha$ . As mentioned in Section 3, these relations are not rigorously proved, even though we know that relations of this form exist. The value of the sum  $\sum_n c_n L_{16}(\alpha^n)$ , which should be a multiple of  $\pi^{16}$ , turns out to be—at least, to 305 digit accuracy—equal to

$$\frac{904068896677319441471456383959007967274858676781011907976799888762570756}{1391630625} \pi^{16}.$$

$n$	$c_n$
1	31646773047454089991074874872115377409615412289257883821314849289011200
2	25091067534991111610431336422823261608598939493205807887933701844011125
3	6426516851004945630792401712640218924732096012292988922326557153689600
4	13013897237290091458885488420735508994896225052675065321587511815005700
5	25345322420083129704294138983088658472477961828569859945452256855687168
6	7654077937866991663920269850717110624730841668799203698754516215186400
7	-21386348881155389104805966211003822515378446012186901064890133970944000
8	-21266457895771812182877802673643067539198868643537409156120825599294550
9	-1608913679173101993450445560427931226772106165751903603469046557081600
10	-3006602534018037060249498048362745795863844063262283069357027575007928
11	-22235894431165401037596833795954124164938433066802035066390519945625600
12	-22388021786025993535779596394330877113951486019249754962484078194515225
13	41163804442546495646577273482500927058102389473207630788145055334400
14	12174993046884795857107894115241919739043395622527097206004912578483400
15	6844294516616620021500451124873691690268431574889793783745454924496896
16	884113568397643613875687455422133057377443632588809535735835348722400



17 24829097569661524994885437517127853543701708402435233939574620268134400  
 18 28498255925832894617958079027165538540096843712257925301476024055353325  
 19 9166754527778720014326220044587296067183679250386045439949761784217600  
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 21 -16245642940492771699066931364054744350799304973969834884361467289600000  
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 24 -1846808352223918020497518383177026788374351085140391933539543804612050  
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 26 11249447059018351361515061841298059010373082845210017340809155226388025  
 27 15574480717724962535771441206565522595335157033372363898543889893785600  
 28 11313602564500264241445701107299008848994067532467357431502793906232400  
 29 8329120904293952846317524093674042573441422235605490061915675232665600  
 30 5881501014664351408528126106754613171638191830768453338752077170311308  
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Henri Cohen, Centre de Recherche en Mathématiques de Bordeaux, Université Bordeaux I, 33405 Talence, France (cohen@alioth.greco-prog.fr)

Leonard Lewin, Department of Electrical and Computer Engineering, University of Colorado at Boulder, Campus Box 425, Boulder, CO 80309-0425

Don Zagier, Max-Planck-Institut für Mathematik, Gottfried-Claren-Str. 26, 5300 Bonn 3, Germany

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